The C-BRAHMS Project

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Abstract

The C-BRAHMS project develops computational methods for content-based retrieval and analysis of music data. A summary of the recent algorithmic and experimental developments of the project is given. The search engine developed by the project is available at http://www.cs.helsinki.fi/group/cbrahms.

1 Introduction

Content-Based Music Retrieval, or CBMR for short, is a research topic studied rather extensively during the last half decade. One of its famous instances is the so-called “query by humming” or WYHIWYG (What You Hum Is What You Get) application. Given a large database of music called the source, the task is to find excerpts in the database that resemble the most (in a musical way) the hummed query pattern.

This paper introduces our CBMR project called C-BRAHMS (Content-Based Retrieval and Analysis of Harmony and other Music Structures) and its output, the C-BRAHMS engine. The project aims at designing and developing efficient methods for computational problems arising from music comparison, analysis, data mining and retrieval. Currently the project has a focus on retrieving polyphonic music in large scale music databases of symbolically encoded music.

The C-BRAHMS project was formally established in January, 2002. C-BRAHMS is part of the From Data to Knowledge (FDK) research unit hosted by the Department of Computer Science at University of Helsinki. FDK has been selected as a centre of excellence funded by the Academy of Finland. The group collaborates with several researchers and research groups abroad.

2 Music retrieval algorithms

Symbolic music data can be seen as strings of symbols, and string-matching-based methods have been applied to CBMR problems. These methods are designed for handling one-dimensional data, and hence they do not apply on polyphonic music without modifications.

The C-BRAHMS project has developed some generalized string matching algorithms that can deal with polyphonic music. For instance the algorithm by Lemström & Tarhio (2003) uses bit-parallelism and precomputed offline data structure containing pitch interval classes for each chord. This structure is then scanned to filter out match candidates, which are checked with another, slower, algorithm.

Our other string-matching-based algorithms allow efficient transposition-invariant approximate searching (Mäkinen, Navarro & Ukkonen, 2003; Lemström & Navarro, 2003); here approximation means insertions and deletions of notes but not small variations of pitch levels. The efficiency is achieved by using sparse dynamic programming and bit-parallel techniques. Similar techniques are used for finding the minimum splitting of a pattern in a multi-track musical work (Lemström & Mäkinen, 2003).

The project has put effort in developing a recent methodology interpreting music data as geometric objects in an Euclidean space (Wiggins, Lemström & Meredith, 2002; Ukkonen, Lemström & Mäkinen, 2003). In geometric representation both the query pattern and the source are represented by objects in a multidimensional space. For instance, in a 2-dimensional time-pitch space, point objects give the onset time and the pitch, while line segment objects (the well-known piano-roll representation) give the duration of the associated notes, as well. To include additional note parameters the dimensionality can be increased without the need to modify algorithms. The approach is inherently transposition-invariant, and dealing with monophonic and polyphonic music is equally straightforward. Moreover, musical decorations, such as ornamentations for instance, do not deteriorate its working.

One of the basic ideas of the geometric approach is to calculate difference vectors between each source point object and pattern point object, sort the vectors, and calculate the frequencies of all such difference vectors. There is a complete match if the frequency of some difference vector equals the number of point objects in the pattern. Then the onset times must match, which means that rhythmic information is taken into account in the geometric approach; note that the string matching approach usually loses this because only the relative order of notes is taken into account. A more efficient version sorts the difference vectors online by using a hash table (Wiggins, Lemström & Meredith, 2002; Ukkonen, Lemström & Mäkinen, 2003).
Small deviations of note onset times inherent in MIDI data generated by playing with a MIDI keyboard confuse the basic geometric approach described above. This problem is solved by another algorithm working on line segment objects, which finds the maximal overlap (common duration) between the line segments of the pattern and the data (Ukkonen, Lemström & Mäkinen, 2003).

To the public demo engine, new algorithms are added after a testing procedure conducted with the private engine. The public demo engine contains most of the algorithms mentioned above.

We provide text-based and WWW based query interfaces. The WWW interface includes a piano keyboard for recording, playing back and refining the query pattern before executing the query. It also allows for choosing the query algorithm, and tuning the error threshold and the number of query results shown to the user. Moreover, the user may select whether the engine gives only the best match or all matches for one database document. The query results are given in a decreasing order of similarity, which may be based on e.g. the amount of transposition and the number of errors in an approximate query task.

The query results include metadata, such as name of the composer; title; opus number; date and genre of the music piece, and content data, such as score files in PostScript and PDF formats; approximate bar number of the match; pitch class names (and octaves) of the matching notes; and the amount of transposition required. The interface allows the matched part or the whole music piece to be played, and to view various histograms of the note data.

The algorithms have been implemented as C language extensions to a distributed server implemented in Ruby language. The mixing of interpreted and compiled code allows for greater productivity in implementing parts that are not essential as regards their performances and their results to similar queries on a given database.

The demo engine uses MIDI files from the Mutopia project. Files are released either as public domain or under a Mutopia-bSD license. Each MIDI file is converted to a string containing variable-length chords which, subsequently, contain notes. This string is included in a object representing the music piece, and searches are performed separately for each piece. According to our experiments, this procedure minimizes the amount of used working memory in contrast to the approach of merging all data into one large string and performing a single search on it (the latter approach would also require removing matches that go beyond music piece boundaries). This is an essential issue with algorithms requiring a large amount of space, when working on large databases and large amounts of concurrent queries are allowed.

According to our experiments, the former arrangement also improves the performance of memory-allocation-intensive algorithms over the latter approach. However, algorithms that do not need memory allocation during search, such as our bit-parallel algorithms, suffer minor performance degradation.

## References


