Sharper Upper Bounds for Unbalanced Uniquely Decodable Code Pairs

Per Austrin, Petteri Kaski, Mikko Koivisto, and Jesper Nederlof

Abstract—Two sets of 0–1 vectors of fixed length form a uniquely decodable code pair if their Cartesian product is of the same size as their sumset, where the addition is pointwise over integers. For the size of the sumset of such a pair, van Tilborg has given an upper bound in the general case. Urbanke and Li, and later Ordentlich and Shayevitz, have given better bounds in the unbalanced case, that is, when either of the two sets is sufficiently large. Improvements to the latter bounds are presented.

Index Terms—Additive combinatorics, binary adder channel, isoperimetric inequality, uniquely decodable code pair, zero-error capacity.

I. INTRODUCTION

A CANONICAL problem in multi-user communication theory is how to coordinate unambiguous communication through a multiple access channel, such that several independent senders can simultaneously send as much information as possible to a single receiver (see, e.g., the book by Cover and Thomas [1, Chapter 15]); this could for example occur when several satellites need to send their data to a single terminal.

Unfortunately, despite vast research in the last decades, even in some of the simplest models the zero-error capacity of such communication channels remains far from clear. An extensively investigated and fundamental example is the two-user binary adder channel (BAC). The zero-error capacity of the BAC is equal to the maximum size of the product of the code sizes of a uniquely decodable code pair (UDCP); a pair $A, B \subseteq \{0, 1\}^n$ such that $|A + B| = |A| \cdot |B|$ where $A + B$ denotes the sumset $\{a + b : a \in A, b \in B\}$, and $a + b$ denotes addition over $\mathbb{Z}^n$.

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Most previous research on UDCPs has focused on constructions. A basic observation is that, if $A_1, B_1 \subseteq 2^{[n]}$ is a UDCP and $A_2, B_2 \subseteq 2^{[n]}$ is a UDCP, then $A_1 \times A_2, B_1 \times B_2$ is also a UDCP; here and henceforth, we freely interchange vectors with sets in the natural way. Therefore, for finding asymptotically good constructions for every $n$, it is sufficient to focus on finite $n$. Letting $\alpha$ and $\beta$ denote respectively $\log_2(|A|)/n$ and $\log_2(|B|)/n$, a natural and popular goal is to find a UDCP maximizing $\alpha + \beta$. The most direct construction is to let $A$ be all strings where the first $\beta n$ coordinates are fixed to 0, and $B$ be all strings which use only the first $\beta n$ coordinates. This yields any pair $(\alpha, \beta)$ with $\alpha + \beta = 1$. The simplest non-trivial construction, $A = \{00, 01, 11\}$, $B = \{10, 01\}$ giving $\alpha + \beta = (\log_2(3) + 1)/2 \approx 1.29248$, was presented by Kasami and Lin [2]. This was the best until 1985. Then it was improved to 1.30366 by van den Braak and van Tilborg [3], and after subsequent improvements by Ahlswede and Balakirsky [4] (1.30369), van den Braak [5] (1.30565), Urbanke and Li [6] (1.30999), the current record is 1.31781 by Mattas and Östergård [7]. Several of these results were obtained by computer searches for finite $n$. More relevant to our study is the important work by Kasami et al. [8], which shows that for sufficiently large $n$ there exist (somewhat surprisingly) UDCPs with $\alpha \geq 1 - o(1)$ and $\beta \geq 0.25$.

Considering upper bounds, the rather direct $\alpha + \beta \leq 1.5$ has been independently found by at least Liao [9], Ahlswede [10], Lindström [11] and van Tilborg [12]. Leaving a gap to the lower bound, 1.5 is still the best upper bound known on $\alpha + \beta$ in general. However, Urbanke and Li [6] managed to break through the 1.5 bound in the unbalanced case: assuming $\alpha \geq 1 - \epsilon$ for a sufficiently small value of $\epsilon$, they showed that $\beta \leq 0.4921$. On a high level, their approach works as follows: a result of van Tilborg [12] (see Lemma 1 below) shows there are not many pairs $(a, b) \in A \times B$ of small Hamming distance, and if $A$ and $B$ are sufficiently large, then the number of such pairs is bounded from below by an isoperimetric inequality for which the authors use Harper’s theorem. Later, this result was improved to $\beta \leq 0.4798$ by Ordentlich and Shayevitz [13]. Their proof idea is somewhat more involved: the authors give a procedure that, given a UDCP $A, B \subseteq \{0, 1\}^n$, constructs another UDCP $C, D \in \{0, 1\}^{(1-\gamma)n}$ of comparable size for some $\gamma > 0$. This was achieved by proving the existence of a subset $L \subseteq [n]$ with $|L| = \gamma n$ such that for some $c \in \{0, 1, 2\}^{|L|}$, the projection $(a + b)_L$ equals $c$ for many pairs $a, b$. The existence of such a subset is proved using a variant of the Sauer–Perles–Shelah lemma. Unfortunately, both the referred bounds [6], [13] converge fast to $(1 - \epsilon) + \beta \leq 1.5$ as $\epsilon$ increases (see Figure 1 of Ordentlich and Shayevitz [13]).
The present authors [14] gave a novel and direct connection between UDCPs and additive number theory. Motivated by algorithm design for the Subset Sum problem, they observed the following: if \( w \in \mathbb{Z}^n, t \in \mathbb{Z} \) and \( A \subseteq \{0, 1\}^n \) such that \( a \cdot w = a' \cdot w \) implies \( a = a' \) for every \( a, a' \in A \), and \( B = \{ b \in \{0, 1\}^n : w \cdot b = t \} \), then \( A, B \) is a UDCP. Here \( \cdot \) denotes the inner product.

The channel capacity application has also inspired studies of several variants of the basic setting of this paper, for example, with both sets being the same [15], [16], with noise [17], or with more than two users [10], [18], [19].

Contributions

Motivated by the lack of progress on the large gap between the current lower and upper bounds for UDCPs, we propose to restrict attention to the case \( |A| \geq 2^{1-\epsilon}n \) for small values of \( \epsilon \): before we can understand the exact tradeoff between \( \alpha \) and \( \beta \), we first need to understand this tradeoff for large values of \( \alpha \). A natural question is whether \( \alpha \geq 1 - o(1) \) implies \( \beta \leq 0.25 + o(1) \); in other words, is the construction of Kasami et al. [8] optimal, or could it be improved? While the present work does not settle this question, we narrow the gap by pushing the upper bound closer to 0.25. Our main result is the following:

**Theorem 1 (Main):** Suppose \( A, B \subseteq \{0, 1\}^n \) is a UDCP with \( |A| \geq 2^{1-\epsilon}n \) and \( |B| = 2^n \). Then \( \beta \leq 0.4228 + \sqrt{\epsilon} \).

Our proof combines ideas from both previous upper bounds [6], [13] with new ideas. We will present our proof by first providing a “warm-up” bound of \( \beta \leq 0.4777 + O(\sqrt{\epsilon}) \) (Theorem 2). To establish this bound, we study the joint probability \( \Pr[a \in A, b \in B] \) for two correlated random vectors \( a, b \in \{0, 1\}^n \). We bound this probability from above and below using, respectively, van Tilborg’s lemma (Lemma 1) and an isoperimetric inequality due to Mossel et al. [20]. This approach is similar to that of Urbanke and Li [6], but improves their bound for small values of \( \epsilon \).

The intuition behind our main bound (and also, in part, the bounds of Urbanke and Li [6] and Ordentlich and Shayevitz [13]) is as follows. The above strategy does not give a good bound if \( A \) and \( B \) are antipodal Hamming balls: the studied probability is very small in this case, so the upper bound is not really stringent. However, intuitively such a pair cannot form a large UDCP since the pairwise sums will be concentrated on the sum of the two centers of the Hamming balls. Our novel approach is that we use the encoding argument from van Tilborg’s lemma to show that if \( A \) is large enough, then \( B \) needs to be sufficiently spread out over the hypercube. Specifically, we show that there exists a set \( L \subseteq [n] \) of size close to \( n/2 \) such that \( L \) has an almost maximum number of projections on \( B \). Subsequently, we use this set \( L \) to define a refined distribution of the vectors \( x \) and \( y \). In the refined distribution, \( x, y \) are only correlated in the coordinates from \( L \), and for applying the isoperimetric inequality the large number of projections is then essential.

II. NOTATION AND PRELIMINARIES

A. Notation

Given reals \( a, b \) with \( b \geq 0 \), we write \( a \pm b \) for the interval \([a-b, a+b]\). If \( n \) is an integer, we denote by \([n]\) the set \( \{1, \ldots, n\} \).

For a vector \( x \in \mathbb{R}^n \), we let \( x^{-1}(z) \subseteq [n] \) denote the set of coordinates \( i \) such that \( x_i = z \). For binary vectors, we apply the usual set operations in the obvious way, by interpreting a vector \( x \in \{0, 1\}^n \) as the set \( x^{-1}(1) \subseteq [n] \).

For example, \( x \setminus y \) is a vector whose \( i \)-th entry is 1 if \( x_i = 1 \) and \( y_i = 0 \), and 0 otherwise; \( x \cap y \) denotes the symmetric difference (or alternatively, the componentwise XOR) of \( x \) and \( y \); and \( |x| \) denotes the Hamming weight of \( x \). Given a vector \( x \in \{0, 1\}^n \) and a subset \( P \subseteq [n] \), we let \( x_P \) denote the projection of \( x \) on \( P \): \( x_P \in \{0, 1\}^P \) such that \( x_P \) agrees with \( x \) on all coordinates in \( P \). For a family \( X \subseteq \{0, 1\}^n \) we also write \( X_P \) for the family \( \{x_P : x \in X\} \).

We write \( o(1) \) for all terms that tend to zero when \( n \) tends to infinity. Such terms can be safely ignored for our purposes as no other variables will depend on \( n \) and upper bounds for UDCPs of large dimension imply upper bounds for UDCPs of finite dimension due to the construction mention in Section I.

B. Entropy

For a real \( x \in [0, 1] \) we denote by \( h(x) \) the binary entropy of \( x \), that is, \( h(x) = -x \log_2(x) - (1-x) \log_2(1-x) \). It is well known that \( h(x) \) is monotone increasing for \( x \in [0, 1/2] \), monotone decreasing for \( x \in [1/2, 1] \), and that \( (h)'' \leq 2h((n)/m) \).

The following inequality can be shown by standard calculus: **Observation 1:** For all \( x \in [0, 1/2] \), \( h(\frac{1}{2} + x) < -\frac{1}{2m^2}x^2 \).

This observation implies another useful bound: **Observation 2:** Let \( \epsilon > 0 \) be a constant. Let \( X \subseteq \{0, 1\}^n \) such that \( |X| \geq 2^{1-\epsilon}n \), \( z \in \{0, 1\}^n \), and \( \gamma \geq \sqrt{\ln(2)/\epsilon^2} \).

Then \( |\{x \in X : |x \triangle z| \in (\frac{1}{2} \pm \gamma)\}| \geq |X|/2 \) for all sufficiently large \( n \).

To see this, note that

\[
|\{x \in X : |x \triangle z| \notin (\frac{1}{2} \pm \gamma)\}| \leq 2 \sum_{k=0}^{\left\lfloor \frac{1}{2\gamma} \right\rfloor} \binom{n}{k} \\
\leq n^2 - \frac{1}{2} - \gamma) \leq n^2 - n^2 \frac{1}{2} - \gamma)
\]

Since \( h(\frac{1}{2} - \gamma) < 1 - \frac{1}{2m^2} \gamma^2 = 1 - \epsilon \), there is some \( \epsilon' > 0 \) (depending only on \( \epsilon \)) such that \( h(\frac{1}{2} - \gamma) = 1 - \epsilon - \epsilon' \). Thus

\[
n^2\gamma(\frac{1}{2} - \gamma) = n^2 - n^2 \epsilon' 2(1-\epsilon) n = n^2 - n^2 \epsilon' |X|
\]

which, for all sufficiently large \( n \), is smaller than \( |X|/2 \).

C. UDCPs

We will use the following well known property of UDCPs that directly follows from noting that whenever \( a - b = a' - b' \) we have \( a + b' = a' + b \).

**Observation 3:** If \( A, B \) is a UDCP, then \( |A - B| = |A| \cdot |B| \).

We will also use the following bound. Since the proof is elegant and highly instructive for understanding our approach, we provide a (known) proof.
Theorem 2: Suppose \( A, B \subseteq \{0,1\}^n \) is a UDCP with \(|A| \geq 2^{(1-\epsilon)n}\) and \(|B| \geq 2^3n\). Then \( \beta \leq 0.4777 + \epsilon + 0.7676\sqrt{\epsilon(1-\beta)} \).

Proof: Let \( W_d = \{(a, b) \in A \times B : |a \triangle b| = d\} \). By definition of \( \rho \)-correlation it is easy to see that
\[
\Pr_{a \sim \rho, b} [a \in A, b \in B] = 2^{-n} \sum_{d=0}^{n} \binom{n}{d} \left( \frac{1}{2} \right)^{n-d} \frac{1-\rho}{2} \left( \frac{1-\rho}{2} \right)^d |W_d| \\
\leq 2^{-2n} \sum_{d=0}^{n} \binom{n}{d} \left( \frac{1+\rho}{2} \right)^{n-d} \left( \frac{1-\rho}{2} \right)^d \\
= 2^{-2n}(3-\rho)^n,
\]
where the inequality follows from Lemma 1, and the last equality follows from the binomial theorem. On the other hand, using Lemma 2, we have that
\[
\Pr_{a \sim \rho, b} [a \in A, b \in B] \geq 2^{-n} \binom{n}{d}(1+\epsilon)\frac{1+\rho}{2}\frac{1-\rho}{2} \sqrt{\epsilon(1-\beta)}
\]
Combining the bounds, taking logs, and dividing by \( n \), we see that for any \( 0 \leq \rho < 1 \),
\[
\frac{1+\beta + 2\sqrt{\epsilon(1-\beta)}}{1-\rho^2} \leq \log_2 (3-\rho) - \frac{2}{\sqrt{n}}
\]
or equivalently,
\[
\beta \leq (\log_2 (3-\rho) - 2)(1-\rho^2) + 1 + 2\rho \sqrt{\epsilon(1-\beta)}.
\]
By setting \( \rho = 0.3838 \) we obtain the claimed bound.

We remark that the proof of Theorem 2 does not use the full strength of Lemma 1. In particular, it only uses that \(|W_d| \leq (\frac{n}{d})^{2d} \). However, using the sharper bound of \((\frac{n}{d})^{2d} \) does not yield any improvement in the exponent because for \( \rho \geq 0 \) the dominating terms in the exponential sum are those where \( d \leq n/2 \).

IV. PROOF OVERVIEW OF MAIN BOUND

The proof of our main bound follows the same blueprint as the proof of Theorem 2, but we use a more refined version of the noise distribution. In particular, we only apply the noise on a subset of \([n]\) where both \( A \) and \( B \) are sufficiently dense, e.g. have sufficiently many projections to that subset.

Definition 1: Fix \( L \subseteq [n] \). Given \( x \in \{0,1\}^n \) let \( y \sim_{\rho} x \) denote that \( y \in \{0,1\}^n \) is the random variable distributed as follows:
\[
y_i \sim_{\rho} x_i, \quad \text{if } i \in L,
\]
\[
y_i \sim_{0} x_i, \quad \text{if } i \notin L.
\]

In other words, \( y \) is a \( \rho \)-correlated copy of \( x \) on the coordinates of \( L \), and uniformly random outside \( L \).

We proceed to give upper and lower bounds on the quantity \( \Pr_{a \sim \rho, b} [a \in A, b \in B] \). In order for these bounds to hold, we need a mild density condition on \( A \) with respect to the split \( (L, [n] \setminus L) \). In particular, we make the following definition.

Definition 2: We say that \( A \subseteq \{0,1\}^n \) is \( \epsilon \)-dense with respect to \( L \subseteq [n] \) if \(|A| \geq 2^{|L|-(1-\epsilon)n}\). Then for any \( L \subseteq [n] \), there is an \( A' \subseteq A \) that is \( \epsilon \)-dense with respect to \( L \).

Proof: For \( a, a' \in A \) note that the condition \( a_L = a'_L \) is an equivalence relation partitioning \( A \) into at most \( 2^{|L|} \) equivalence classes, each of size at most \( 2^{n-|L|} \). It follows that there must be at least \(|A|/2^{n-|L|} \geq 2^{|L|-(1-\epsilon)n} \) equivalence classes of size at least \(|A|/2^{|L|+1} = 2^{n-|L|-(1-\epsilon)n} \) and we can take \( A' \) to be the union of these.

With these definitions in place, we are ready to state the precise upper and lower bounds on the refined noise probability.
Lemma 3: Fix $L \subseteq [n]$ and let $\lambda = |L|/n$. Then for any $0 \leq \rho \leq 1$ and UDCPs $(A, B)$ such that $A$ is $\epsilon$-dense with respect to $L$, we have

$$\log_2 \Pr_{u \sim \rho L}[a \in A, b \in B] \leq \frac{\sqrt{\ln(2)e^{-\frac{\lambda}{2}}} - \frac{1}{2}}{n} + \lambda \left(\log_2(3 - \rho) - \frac{3 \rho}{2}\right) + o(1).$$

The proof appears in Section VI.

Lemma 4: Fix $L \subseteq [n]$ and let $\lambda = |L|/n$. Then for any $0 \leq \rho < 1$ and UDCPs $(A, B)$ such that $A$ is $\epsilon$-dense with respect to $L$ and $|B_L| = 2^{\rho n}$ for some $0 \leq \pi \leq \lambda$, we have

$$\log_2 \Pr_{u \sim \rho \lambda L}[a \in A, b \in B] \geq \frac{\pi - \lambda - \epsilon - 2\rho \sqrt{\ln(2)e^{-\frac{\lambda}{2}}}}{n} - \frac{1}{2} \pi^2 + \lambda - 1 - \epsilon - o(1).$$

The proof appears in Section VII.

The proof of Lemma 4 may be found in Section VII.

Lemma 5: For sufficiently large $n$ and UDCPs $(A, B)$ such that $|A| \geq 2^{(1-\epsilon)n}$, $|B| = 2^{\rho n}$, there exists $L \subseteq [n]$ such that $|L| \leq \frac{1}{2} \ln(2)e^{-\frac{\lambda}{2}}$ and $|B_L| \geq 2^{(\beta - \epsilon)n-1}$.

Proof: Let $F \subseteq A \times B$ consist of all pairs $(a, b)$ such that $|a \Delta b| \in (\frac{1}{2} \pm \ln(2)e^{-\frac{\lambda}{2}})n$. We have that

$$|P| = \sum_{b \in B} \left| \{a \in A : |a \Delta b| \in (\frac{1}{2} \pm \ln(2)e^{-\frac{\lambda}{2}})n \} \right| \geq \sum_{b \in B} |A|/2 = |A| \cdot |B|/2,$$

where the inequality is by Observation 2. Similarly as in the proof of Lemma 1, consider the encoding

$$h : (a, b) \mapsto (a \Delta b, b \setminus a).$$

By Observation 3, $|A - B| = |A| \cdot |B|$, and since $a - b$ can be computed from $h(a, b)$, it follows that $h$ is injective and $|h(P)| = |P|$. We now bound $|h(P)|$ from above. To this end, note that $b \setminus a \subseteq a \Delta b$, and so $b \setminus a \in B_{a \Delta b}$. (More precisely, $b \setminus a$ projected to $a \Delta b$ is in $B_{a \Delta b}$; we only need that $b \setminus a$ can be described by a single element of $B_{a \Delta b}$.) Therefore, by summing over the possible values of $X = a \Delta b$ we have that

$$|h(P)| \leq \sum_{X \subseteq [n]} |B_X|.$$

Thus there must be an $X \subseteq [n]$ with $|X| \in (\frac{1}{2} \pm \ln(2)e^{-\frac{\lambda}{2}})n$ and $|B_X| \geq |h(P)|/2^n = |P|/2^n \geq |A| \cdot |B|/2^{n+1} \geq 2^{(\beta - \epsilon)n-1}$, as we claimed.

VI. UPPER BOUND: PROOF OF LEMMA 3

In this section, we prove the upper bound on the refined noise probability stated in Lemma 3. Fix $L \subseteq [n]$ and let $\lambda = |L|/n$. Furthermore, let $0 \leq \rho \leq 1$ and let $(A, B)$ be a UDCP such that $|A|$ is $\epsilon$-dense with respect to $L$.

Let $R = |L| \setminus L$ be the coordinates not in $L$. Let $W_d$ be the set of pairs $a_{L \Delta R} \in A, b_{L \Delta R} \in B$ such that $|a_{L \Delta R}| = d$. Theorem 2: For sufficiently large $n$, we have that

$$|W_d| \leq \binom{|L|}{d} 2^{d^2 2^L |R| 2^n \ln(2)e^{-\frac{\lambda}{2}}/2 + 2}.$$
Proof: Let \( \epsilon' = \sqrt{(\epsilon \ln 2)/(2(1 - \lambda))} \), and let \( W'_d \subseteq W_d \) be all pairs from \( W_d \) such that \( \frac{|a_R \triangle b_R|}{|R|} \in \frac{1}{2} \pm \epsilon' \). Similarly as in the proof of Lemma 5, we see that

\[
|W'_d| = \sum_{b_L,b_R \in B \atop a_L \in A_L \atop |a_L \triangle b_R| = d} \left| \{ a_R \in \{ 0, 1 \}^R : |a_R \triangle b_R| \in (\frac{1}{2} \pm \epsilon') |R| \} \right|
\]

\[
\geq \sum_{b_L,b_R \in B \atop a_L \in A_L \atop |a_L \triangle b_R| = d} \frac{1}{2} |\{ a_R \in \{ 0, 1 \}^R : a_R \in A \} | = \frac{1}{2} |W_d| .
\]

The inequality follows from Observation 2 combined with the \( \epsilon' \)-dense property:

\[
|\{ a_R \in \{ 0, 1 \}^R : a_R \in A \} | \geq 2^{|R| - \epsilon n - 1} = 2^{(1 - \frac{\epsilon}{\ln 2}) |R| - 1} .
\]

We proceed to bound \( |W'_d| \) from above. Similarly as in the proof of Lemma 1, define an encoding \( \eta \) on elements \( (a, b) \) of \( W'_d \):

\[
\eta : (a_L a_R, b_L b_R) \mapsto (a_L \triangle b_L, a_L \setminus b_L, a_L \triangle b_R, a_R \setminus b_R).
\]

Since the image \( \eta(a, b) \) directly gives \( a - b \) and \( |A| - B = |A| |B| \) by Observation 3, we have that \( \eta \) is injective and thus

\[
|W'_d| = |\eta(W'_d)| \leq \left( \frac{|L|}{d} \right)^{2d} \sum_{i \in (0.5 \pm \epsilon')} \left( \frac{|R|}{i} \right)^{2i},
\]

where the inequality follows by bounding the number of possibilities in every coordinate of \( \eta(\cdot) \). The claim is then implied for sufficiently large \( n \) from the easy observation that

\[
\sum_{i \in (0.5 \pm \epsilon')} \left( \frac{|R|}{i} \right)^{2i} \leq 2^{1.5 \epsilon' |R| |R|} \leq 2^{1.5 |R| n + \sqrt{\ln(2) \epsilon}/2} .
\]

By the refined definition of \( \sim'_\rho \) we have that

\[
\Pr_{a \sim'_{\rho} b}[a \in A, b \in B] = 2^{-n} \sum_{d=0}^{|L|} \left( \frac{1 + \rho}{2} \right)^{|L| - d} \left( \frac{1 - \rho}{2} \right)^{d} 2^{-|R| |W_d|} . (4)
\]

To see this, note that \( |W_d| \) counts exactly the pairs \( a \in A, b \in B \) satisfying \( |a_L \triangle b_L| = d \), and that the probability that such a pair is picked can be computed as the probability that \( a \) is picked (which is \( 2^{-n} \)) times the probability that \( b \) is picked given that \( a \) is given. The probability that \( b_R \) is picked is simply \( 2^{-|R|} \) since it is picked uniformly at random, and the probability that \( b_L \) is picked is \( \left( \frac{1 + \rho}{2} \right)^{|L| - d} \left( \frac{1 - \rho}{2} \right)^{d} \), similarly as in the proof of Theorem 2.

Using Claim 2, we bound (4) from above by

\[
\Pr_{a \sim'_{\rho} b}[a \in A, b \in B] \leq 2^{-n} \sum_{d=0}^{|L|} \left( 1 + \rho \right)^{|L| - d} \left( 1 - \rho \right)^{d} \left( \frac{|L|}{d} \right)^{2d} 2^{1.5 |R| n + \sqrt{\ln(2) \epsilon}/2 + 1}
\]

\[
= 2^{-n + 1.5 |R| n + \sqrt{\ln(2) \epsilon}/2 + 1} \sum_{d=0}^{|L|} \left( 1 + \rho \right)^{|L| - d} \left( 2 - 2 \rho \right)^{d} \left( \frac{|L|}{d} \right)^{2d}
\]

\[
= 2^{\sqrt{\ln(2) \epsilon}/2 - 2} n^{1.5 |R| + 1} (3 - \rho)^{1} |L| ,
\]

where the last equality follows from the binomial theorem. Using \( |R| = n - |L| \), taking logs, and dividing by \( n \), we get

\[
\log_2 \Pr_{a \sim'_{\rho} b}[a \in A, b \in B] \leq \sqrt{\ln(2) \epsilon} - \frac{1}{2} + \lambda \left( \log_2 (3 - \rho) - \frac{3}{2} \right) + 1/n.
\]

VII. LOWER BOUND: PROOF OF LEMMA 4

In this section, we prove the lower bound on the refined noise probability stated in Lemma 4. Fix \( L \subseteq \{ n \} \) and let \( \lambda = |L|/n \). Furthermore, let \( 0 \leq \rho < 1 \), and let \( (A, B) \) be a UDCP such that \( A \) is \( \epsilon \)-dense with respect to \( L \) and \( |B_L| = 2^{|n|} \) for some \( 0 \leq \pi \leq \lambda \).

Due to the chain rule

\[
\Pr_{a \sim'_{\rho} b}[a \in A, b \in B] = \Pr_{a \sim'_{\rho} b}[a \in A, b \in B\mid a_L \in A_L, b_L \in B_L] \times \Pr_{a \sim'_{\rho} b}[a_L \in A_L, b_L \in B_L]. (5)
\]

We proceed by giving lower bounds for the two factors in the product (5). Let \( R = \{ n \} \setminus L \). For the first factor, note that if \( b_L \in B_L \), there is at least one \( b_R \) such that \( b_L b_R \in B \) by the definition of \( B_L \), and such a \( b_R \) is picked with probability \( 2^{-|R|} \) since it is uniformly distributed over \( 2^R \). Similarly, if \( a_L \in A_L \), there are at least \( 2^{|L| - \epsilon n}/2 \) sets \( a_R \subseteq L \) such that \( a_L a_R \in A' \) by the definition of \( A' \), and so such an \( a_R \) is picked with probability at least \( 2^{-\epsilon n}/2 \). In summary,

\[
\Pr_{a \sim'_{\rho} b}[a \in A, b \in B\mid a_L \in A_L, b_L \in B_L] \geq 2^{-|L| - \epsilon n/2} = 2^{(\lambda - 1 - \epsilon - o(1)) n} .
\]

For the second term, apply Theorem 2 with \( U = L \) and

\[
F = A_L, \quad f = \frac{|L| - \epsilon n - 1}{|L|} = 1 - \frac{\epsilon}{\lambda} - o(1),
\]

\[
G = B_L, \quad g = \frac{\pi}{\lambda},
\]

which gives that

\[
\log_2 \Pr_{a \sim'_{\rho} b}[a_L \in A_L, b_L \in B_L] \geq -||L|| \left( \frac{1}{\lambda} + \frac{\pi}{\lambda} + o(1) + 2 \rho \sqrt{\log(1 - \frac{1}{\lambda}) (1 + \pi o(1))} \right)
\]

\[
= n \left( \pi - \lambda - \epsilon - 2 \rho \sqrt{e \lambda - e \pi} - o(1) \right).
\]

The statement now follows by multiplying the two lower bounds into a lower bound for the product (5).

VIII. CONCLUSION

We presented a new upper bound for UDCPs, considerably strengthening previous bounds. We obtained the bound by combining an isoperimetric inequality, which was not used before in the UDCP literature, with an extension of van Tilborg’s bound that works well if the set families are clustered.

Two outstanding open questions that are of main interest remain. In our setting (\( \alpha \geq 1 - \epsilon \)), there is still a big gap between the best construction (achieving \( \beta \geq 1/4 \)) and our new upper bound of \( 0.4228 + \sqrt{\epsilon} \). Narrowing this gap from either direction would be very interesting. In the general case, a major unresolved problem is whether the classic upper bound of \( |A| \cdot |B| \leq 2^{1.5 n} \) is tight.
REFERENCES


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