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Finite-State Spell-Checking with Weighted Language and Error Models—Building and Evaluating Spell-Checkers with Wikipedia as Corpus

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Abstract

In this paper we present simple methods for construction and evaluation of finite-state spell-checking tools using an existing finite-state lexical automaton, freely available finite-state tools and Internet corpora acquired from projects such as Wikipedia. As an example, we use a freely available open-source implementation of Finnish morphology, made with traditional finite-state morphology tools, and demonstrate rapid building of Northern Sámi and English spell checkers from tools and resources available from the Internet.

1. Introduction

Spell-checking is perhaps one of the oldest most researched application in the field of language technology, starting from the mid 20th century (Damerau, 1964). The task of spell-checking can be divided into two categories: isolated non-word errors and context-based real-word errors (Kukich, 1992). This paper concentrates on checking and correcting the first form, but the methods introduced are extendible to context-aware spell-checking.

To check whether a word is spelled correctly, a language model is needed. For this article, we consider a language model to be a one-tape finite-state automaton recognising valid word forms of a language. In many languages, this can be as simple as a word list compiled into a suffix tree automaton. However, for languages with productive morphological processes in compounding and derivation that are capable of creating infinite dictionaries, such as Finnish, a cyclic automaton is required. In order to suggest corrections, the correction algorithm must allow search from an infinite space. A nearest match search from a finite-state automaton is typically required (Oflazer, 1996). The reason we stress this limitation caused by morphologically complex languages is that often even recent methods for optimizing speed or accuracy suggest that we can rely on finite dictionaries or acyclic finite automata as language models.

To generate correctly spelled words from a misspelled word form, an error model is needed. The most traditional and common error model is the Levenshtein edit distance, attributed to (Levenshtein, 1966). In the edit distance algorithm, the misspelling is assumed to be a finite number of operations applied to characters of a string: deletion, insertion, change, or transposition1. The field of approximate string matching has been extensively studied since the mid 20th century, yielding efficient algorithms for simple string-to-string correction. For a good survey, see (Kukich, 1992). Research on approximate string matching has also provided different fuzzy search algorithms for finding the nearest match in a finite-state representation of dictionaries.

For the purpose of the article, we consider the error model to be any two-tape finite state automaton mapping any string of the error model alphabet to at least one string of the language model alphabet. As an actual implementation of Finnish spell-checking, we use a finite-state implementation of a traditional edit distance algorithm. In the literature, the edit distance model has usually been found to cover over 80% of the misspellings at distance one (Damerau, 1964). Furthermore, as Finnish has a more or less phonemically motivated orthography, the existence of homophonic misspellings are virtually non-existent. In other words, our base assumption is that biggest source of errors for Finnish spell checking is slip of finger style of typos, for which the edit distance is a good error model.

The statistical foundation for the language model and the error model in this article is similar to the one described in (Norvig, 2010), which also gives a good overview of the statistical basis for the spelling error correction problem along with a simple and usable python implementation.

For practical applications, the spell-checker typically needs to provide a small selection of the best matches for the user to select from in a relatively short time span, which means that when defining corrections, it is also necessary to specify their likelihood in order to rank the correction suggestions. In this article, we show how to use a standard weighted finite-state framework to include probability estimates for both the language model and the error model. For the language model, we use simple unigram training with a Wikipedia corpus with the more common word forms to be suggested before the less common word forms. In the error model, we design the weights in the edit distance automaton so that suggestions with a greater Levenshtein-Damerau edit distance are suggested after those with fewer errors.

To evaluate the spell-checker even in the simple case of correcting non-word errors in isolation, a corpus of spelling mistakes with expected corrections is needed. A construction of such a corpus typically requires some amount of manual labour. In this paper, we evaluate the test results both against a manually collected misspelling corpus and against automatically misspelled texts. For a description of the error generation techniques, see (Bigert, 2005).

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1 Transposition is often attributed to an extended Levenshtein-Damerau edit distance given in (Damerau, 1964)
2. Goal of the paper

In this article, we demonstrate how to build and evaluate a spell-checking and correction functionality from an existing lexical automaton. We present a simple way to use an arbitrary string-to-string relation transducer as a misspelling model for the correction suggestion algorithm, and test it by implementing a finite-state form of the Levenshtein-Damerau edit distance relation. We also present a unigram training method to automatically rank spelling corrections, and evaluate the improvement our method brings over a correction algorithm using only the edit distance. The paper describes a work-in-progress version of a finite state spell-checking method with instructions for building the spellers for various languages and from various resources. The language model in the article is an existing free open-source implementation of Finnish morphology\(^2\) (Pirinen, 2008) compiled with HFST (Lindén et al., 2009)—a free, featurewise fully compliant implementation of the traditional Xerox-style LexC and TwolC tools\(^3\). One aim of this paper is to demonstrate the use of Wikipedia as a freely available open-source corpus\(^4\). The Wikipedia data is used in this experiment for training the lexical automaton with word form frequencies, as well as collecting a corpus of spelling errors with actual corrections.

The field of spell-checking is already a widely researched topic, cf. the surveys in (Kukich, 1992; Schulz and Mihov, 2002). This article demonstrates a generic way to use freely available resources for building finite-state spell-checkers. The purpose of using a basic finite-state algebra to create spell-checkers in this article is two-fold. Firstly, the amount of commonly known implementations of morphological language models under different finite-state frameworks suggest that a finite-state morphology is feasible as a language model for morphologically complex languages. Secondly, by demonstrating the building of an application for spell-checking with a freely available open-source weighted finite-state library, we hope to outline a generally useful approach to building open-source spell-checkers. To demonstrate the feasibility of building a spell-checker from freely available resources, we use basic composition and n-best-path search with weighted finite-state automata, which allows us to use multiple arbitrary language and error models as permitted by finite-state algebra. To the best of our knowledge, no previous research has used or documented this approach.

To further evaluate plausibility of rapid conversion from morphological or lexical automata to spell checkers we also sought and picked up a free open implementation of Northern Sámi morphological analyzer\(^5\), a word list of English from (Norvig, 2010), and briefly tested them with the same methods and error model as for Finnish. While the main focus of the article is creation and evaluation of Finnish finite-state spell checker, we show examples of building and evaluation for other languages.

3. Methods

The framework for implementing the spell-checking functionality in this article is the finite-state library HFST (Lindén et al., 2009). This requires that the underlying morphological description for spell-checking is compiled into a finite-state automaton. For our Finnish and Northern Sámi examples, we use a traditional linguistic description based on the Xerox LexC/TwolC formalism (Beesley and Karttunen, 2003) to create a lexical transducer that works as a morphological analyzer. As the morphological analyses are not used for the probability weight estimation in this article, the analysis level is merely discarded to get a one-tape automaton serving as a language model. The word list of English is merely compiled to a one tape suffix tree automaton.

The language model can be as simple as a list of words compiled into a suffix tree automaton or as elaborate as a full-fledged morphological description in a finite-state programming language, such as Xerox LexC and TwolC.\(^6\) The words that are found in the transducer are considered correct. The rest are considered misspelled.

It has previously been demonstrated how to add weights to a cyclic finite-state morphology using information on base-form frequencies. The technique is further described in (Lindén and Pirinen, 2009). In the current article, the word form counts are based on data from Wikipedia. The training is in principle a matter of collecting the corpus strings and their frequencies and composing them with the finite-state lexical data. Deviating from the article (Lindén and Pirinen, 2009), we only count full word forms. No provisions for compounding of word forms based on the training data are made, i.e. the training data is composed with the lexical model. This gives us an acyclic lexicon with the frequency data for correctly spelled words.

The actual implementation goes as follows. Clean up the Wikipedia dump to extract the article content from XML and Wikipedia mark-up by removing the mark-up and contents of mark-up that does not constitute running text, leaving only the article content untouched. The tokenization is done by splitting text from white space characters and separating word final punctuation. Next we use the spell-checking automaton to acquire the correctly spelled word forms from the corpora, and count their frequencies. The formula for converting the frequencies \(f\) of a token in the corpus to a weight in the finite-state lexical transducer is

\[
W_{f} = -\log \frac{f}{CS},
\]

where \(CS\) is the corpus size in tokens. The resulting strings with weights can then be compiled into paths of a weighted automaton, i.e. into an acyclic tree automaton with log probability weights in the final states of the word forms. The original language model is then weighted by setting all word forms not found in the corpus to a weight greater than the word with frequency of one, e.g. \(W_{\text{max}} = -\log \frac{1}{CS+1}\). The simplest way to achieve this is to compose the \(\Sigma^+\) automaton with final weight \(W_{\text{max}}\) with

\(^{2}\)http://home.gna.org/omorf1
\(^{3}\)http://hfst.sf.net
\(^{4}\)Database dumps available at http://download.wikimedia.org
\(^{5}\)http://divvun.no
\(^{6}\)It is also possible to convert aspell and hunspell style descriptions into transducers. Preliminary scripts exist in http://hfst.sf.net/.
the unweighted cyclic language model. Finally, we take the union of the cyclic model and the acyclic model. The word forms seen in the corpus will now have two weights, but the lexicon can be pruned to retain only the most likely reading for each string.

For example in Finnish Wikipedia there were 17,479,297 running tokens\(^7\), and the most popular of these is ‘ja’ and with 577,081 tokens, so in this language model the \( W_{ja} = - \log_{10} \frac{17479297}{577081} \approx 4.44 \). The training material is summarized in the Table 1. In the token count is total number of tokens after the preprocessing and tokenization. The unique strings is the number of unique tokens that belonged to the language model, that is the size of actual training data, after unification and discarding potential misspellings and other strings not recognized by language model. For this reason the English training model is rather small, despite the relative size of the corpus, since the finite language model only covered a very small part of unique tokens.

<table>
<thead>
<tr>
<th>Language</th>
<th>Finnish</th>
<th>Northern Sami</th>
<th>English</th>
</tr>
</thead>
<tbody>
<tr>
<td>Token count</td>
<td>17,479,297</td>
<td>258,366</td>
<td>2,110,728,338</td>
</tr>
<tr>
<td>Unique language strings</td>
<td>968,996</td>
<td>44,976</td>
<td>34,920</td>
</tr>
<tr>
<td>Download size</td>
<td>956 MiB</td>
<td>8.7 MiB</td>
<td>5.6 GiB</td>
</tr>
<tr>
<td>Version used</td>
<td>2009-11-17</td>
<td>2010-02-22</td>
<td>2010-01-30</td>
</tr>
</tbody>
</table>

Table 1: Token counts for wikipedia based training material

For finding corrections using the finite-state methodology, multiple approaches with specialized algorithms have been suggested, e.g. (Oflazer, 1996; Schulz and Mihov, 2002; Huldén, 2009). In this article, we use a regular weighted finite-state transducer to represent a mapping of misspellings to correct forms. This allows us to use any weighted finite-state library that implements composition. One of the simplest forms of mapping misspellings to correct strings is the edit distance algorithm usually attributed to (Levenshtein, 1966) and furthermore in the case of spell-checking to (Damerau, 1964). A finite-state automaton representation is given in e.g. (Schulz and Mihov, 2002). A transducer that corrects strings can be any arbitrary string-to-string mapping automaton, and can be weighted. In this article, we build the edit distance mapping transducer allowing two edits. Since the error model can also be weighted, we have applied a weight \( W_{max} \), which is greater than any of the weights given by the language model. As a consequence, our weighted edit distance will function like the traditional edit distance algorithm when generating the corrections for a language model, i.e. any correct string with edit distance one is considered to be a better correction than a misspelling with edit distance two. For example assuming misspelling ‘jq’ for ‘ja’, the error model would find ‘ja’ at distance of \( W_{max} \), but also e.g. ‘jo’ already and so on. In this case the frequency data obtained from Wikipedia will give us the popularity order of ‘ja’ > ‘jo’. A fraction of this weighted edit distance two transducer is given in Figure 1. The transducer in the figure gives full edit distance 2 transducer for language of two alphabet; a transducer for full alphabet is just a union of transducer like that for each pair of alphabets in the language\(^8\).

To get a ranked set of spelling correction suggestions, we simply compile the misspelled word into a path automaton \( T_{word} \). The path automaton is composed with the correction relation \( T_E \)—in this case the weighted edit distance two transducer—to get an automaton that contains all the possible spelling corrections \( T_{sug} = T_{word} \circ T_E \). We then compose the resulting automaton with the original weighted lexical data \( T_L \) to find which string corrections are real words of the language model \( T_f = T_{sug} \circ T_L \). The resulting transducer now contains a union of words along with the combination of weights for the frequency of the word form and the weight for the edit distance. From this transducer, the ranked list of spelling suggestions is gained by a standard n-best-path algorithm listing unique suggestions.

4. Test Data Sets

For the Finnish test material, we use two types of samples extracted from Wikipedia. First, we use a hand-picked selection of 761 misspelled strings found by browsing the strings that the speller rejected. These strings were then manually corrected using a native reader’s best judgement from reading the misspelled word in context to achieve a gold standard for evaluation. Another larger set of approximately 10,000 evaluation strings was created by using the strings from the same

\(^7\)with relatively naive preprocessing and tokenization, splitting at spaces and filtering html and Wikipedia markup

\(^8\)for source code of the Finnish edit distance in HFST framework, see http://svn.gna.org/viewcvs/omorfi/trunk/src/suggestion/edit-distance-2.text
Wikipedia corpus, and automatically introducing spelling errors similar to the approach described in (Bigert et al., 2003), using isolated word Damerau-Levenshtein type errors with a probability of approximately 0.33 \% per character. This error model could also be considered an error model applied in reverse compared to the error model used when correcting misspelled strings. As there is nothing limiting the number of errors generated per word than the word length, this error model may introduce words with an edit distance greater than two.

For Northern Sámi gold standard evaluation we used the test suite included in the svn distribution\(^9\), it seems to be set of common typos of a sort.

For English gold standard evaluation material we use the same Birkbeck spelling error corpus as in article (Norvig, 2010), although the corpus has only restricted free licensing, restrictions being such that it cannot be used e.g. in free open source project.

5. Evaluation

To evaluate the correction algorithm, we use the two data sets introduced in the previous section. However, we use a slightly different error model to automatically correct misspellings than we use for generating them, i.e. some errors exceeding the edit distance of two are unfixable by the error model we use for correction.

The evaluation of the correction suggestion quality is given in Tables 2 and 3. The Table 2 contains precision values for the spelling errors from real texts, and Table 3 for the automatically introduced spelling errors. The precision is measured by ranked suggestions. In the tables, we give results separately for ranks 1—4, and for the remaining lower ranks. The lower ranks ranged from 5—440 where the number of total suggestions ranged from 1—600. In the last column, we have the cases where a correctly written word could not be found with the proposed suggestion algorithm. The tables contain both the results for the weighted edit distance relation, and for a combination of the weighted edit distance relation and the word form frequency data from Wikipedia.\(^10\)

<table>
<thead>
<tr>
<th>Material</th>
<th>Rank 1</th>
<th>2</th>
<th>3</th>
<th>4</th>
<th>Lower</th>
<th>No rank</th>
<th>Total</th>
</tr>
</thead>
<tbody>
<tr>
<td>Finnish</td>
<td>4444</td>
<td>1138</td>
<td>305</td>
<td>301</td>
<td>495</td>
<td>1403</td>
<td>1469</td>
</tr>
<tr>
<td></td>
<td>44 %</td>
<td>11 %</td>
<td>6 %</td>
<td>4 %</td>
<td>11 %</td>
<td>15 %</td>
<td>15 %</td>
</tr>
<tr>
<td>Northern Sámi</td>
<td>1249</td>
<td>267</td>
<td>156</td>
<td>80</td>
<td>828</td>
<td>778</td>
<td>100%</td>
</tr>
<tr>
<td></td>
<td>13 %</td>
<td>3 %</td>
<td>1 %</td>
<td>1 %</td>
<td>5 %</td>
<td>77 %</td>
<td>100%</td>
</tr>
<tr>
<td>English</td>
<td>4422</td>
<td>938</td>
<td>337</td>
<td>290</td>
<td>1353</td>
<td>2657</td>
<td>100%</td>
</tr>
<tr>
<td></td>
<td>44 %</td>
<td>10 %</td>
<td>3 %</td>
<td>3 %</td>
<td>14 %</td>
<td>27 %</td>
<td>100%</td>
</tr>
</tbody>
</table>

Table 3: Precision of suggestion algorithms with generated spelling errors

As a first impression we note that mere Wikipedia training does improve the results in all cases; the number of suggestions in first position rises in all test sets and languages. This would suggest in fact that more mistakes are made in more common words than rare ones, since the low ranking word count did not increase as a result of Wikipedia training.

In Finnish tests, for the actual errors in the real texts’ spelling-error corpus which dominate the lowest ranks of correction suggestions, hapological cases like ‘kokonais-malminvarioista’ from total ore resources spelled as ‘konais-malminvarioista’ came in at the top of the list for both methods, because the correct word form is probably non-existent in the training corpus, and the multi-part productive compound with an ambiguous segmentation produces lots of nearer matches at edit distance one. A more elaborate error model considering hapology as misspelling with a weight equal or less than a single traditional edit distance would of course improve the suggestion quality in this case.

The number of words getting no ranks is common to all methods. They indicate the spelling errors for which the correct form was not among the ones covered by the error model of edit distance two. A good number of these are cases which were not considered in the error model, e.g. a missing space causing run-on words (‘ensisijassa’ instead of ‘ensi sijassa’ in the first place). A good number of mistakes also comes from the use of spoken or informal language forms for very common words, which tend to deviate more than edit distance two (‘esmeks’ instead of ‘esimerkiksi’ for example), with a few more due to missing forms in the language model. E.g. ‘bakterisidin’ is one edit from ‘bakterisidin’ as bactericide, but the correction is not made because the word does not exist in the language model. Both of these error types are correctable by adding words to the lexicon, i.e. the language model, e.g. using special-purpose dictionaries, such as spoken language or medical dictionaries. Finally there is a handful of errors that seem legitimate spelling mistakes of more than two edits (‘assosioitten’ instead of ‘assosiaatioiden’). For these cases, a different error model than the basic edit distance might be necessary.

For Northern Sámi spelling error corpus we notice, that large amount of errors is not covered by the error model. This means that the error model is not sufficient for North-

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\(^{9}\)https://victorio.uit.no/langtech/trunk/gt/sme/src/typos.txt

\(^{10}\)For full tables and test logs, see http://home.gna.org/omorfi/testlogs.
ern Sámi spell checking, as we can see a number of errors with edit distance of greater than 2 (e.g. ‘sáđdejuvvon’ instead of ‘sáđdejeuvvon’).

Comparing our English test results with previous research using programmatic implementation of same language and error model, we first note again that great number of words are out of reach by error model of mere edit distance 2, some of them are even real word spelling errors, such as ‘gone’ in stead of ‘went’, but unfortunately they were intermixed with other spelling error material so we did not have time to weed them out from corpus. The rest of spelling errors beyond edit distance 2 are mostly caused by English orthography being relatively distant from pronunciation, such as ‘negoshayshauns’ in stead of ‘negotiations’, which usually are corrected with very different error models such as soundex or other phonetic keys as demonstrated in e.g. (Mitton, 2009). The results of evaluation of correction suggestions for the testing materials show similar tendency as found in the original article (Norvig, 2010).

The impact on performance of using non-optimized methods to check spelling and get suggestion lists was not thoroughly measured, but to give an impression of methods general applicability, we note that for Finnish material of generated misspellings, the speed of spell-checking was 3.18 seconds for 10,000 words or approx. 3,000 words per second, and the speed of generating suggestion lists (all possible corrections) for misspelled words took 10,493 seconds for 10,000 words or approx. 1 word per second, when measured using GNU time utility and hfst-omor-evaluate program from HFST toolset, which batch processes spell-checking tasks on tokenized input and evaluates precision and recall against correction corpus. The space requirements for Finnish spell checking automata are 9.2 MiB for Finnish morphology and 378 KiB for Finnish edit distance 2 automaton for alphabet of size 72. As a comparison the English language model obtained from the word list is only 3.2 MiB in size, and correspondingly the error model 273 KiB for alphabet of size 54.

6. Discussion

The obvious and popular development is to extend the language model to support n-gram dictionaries with $n > 1$, which has been shown to be a successful technique for English e.g. in (Mays et al., 1991). The extension using the same framework is not altogether trivial for misspelled words took 10,493 seconds for 10,000 words or approx. 1 word per second, when measured using GNU time utility and hfst-omor-evaluate program from HFST toolset, which batch processes spell-checking tasks on tokenized input and evaluates precision and recall against correction corpus. The space requirements for Finnish spell checking automata are 9.2 MiB for Finnish morphology and 378 KiB for Finnish edit distance 2 automaton for alphabet of size 72. As a comparison the English language model obtained from the word list is only 3.2 MiB in size, and correspondingly the error model 273 KiB for alphabet of size 54.

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7. Conclusions

In this article we have demonstrated an approach for creating spell-checkers from various language models—ranging from as simple as word list to as complex as full fledged implementation of morphology—built into a finite-state automata. We also demonstrated simple approach to training the models using word frequency data extracted from Wikipedia. Further, we have presented a construction of a simple edit distance error model in the form of a weighted finite-state transducer, and proven usability of this basic finite state approach by evaluating produced spell-checkers against both manually collected smaller and automatically created larger error corpora. Given the amount of finite-state implementations of morphological language models it seems reasonable to expect support for use of general finite-state methodology and language models is possible to support spell-checking for large array of languages. For the fact that arbitrary weighted two-tape automaton may be used to implement error model, it will also be possible to easily extend the models with basic available finite-state toolsets without developing additional tools. We also showed that combining the basic edit distance error model with a simple unigram frequency model already improves the quality of the error corrections. We also note that even using basic finite-state transducer algebra from freely available finite-state toolkits and no specialized algorithms, the speed and memory requirements of the spell-checking seems sufficient for typical interactive usage.

8. Acknowledgements

We thank the colleagues in HFST research team and anonymous reviewers for valuable suggestions regarding the article.

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11 http://foma.sf.net

12 http://hunspell.sf.net

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